

# COMP4019 - Lab Session 7 – Binary Heaps

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## 1 Implicit Data Structures

Consider the representation of a *complete  $m$ -ary tree* (that is, a tree where each node has  $k$  children, and such that all levels are filled, except possibly the last one, which is flushed to the left). See Figure 1 for an example. We have discussed how this type of tree can be represented *implicitly* in an array. For this, we need to map each node in the tree to an index in the array.

1. For a tree of a given branching factor  $m$  and a maximum depth  $d$ , what size should the array be initialized to?
2. Given a list  $(x_1, \dots, x_k)$ , where  $0 \leq x_i < m$  and  $k \leq d$  representing a path to a node in the tree, find the corresponding index in the array;
3. What about the inverse operation?

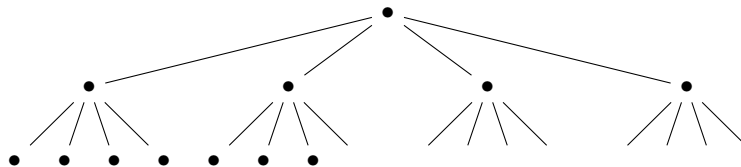


Figure 1: A complete 4-ary tree.

## 2 Binary Heaps

Implement a binary heap in your favorite programming language (most specifically the operations `insert`, `minimum`, and `extract`). Depending on the programming paradigm you choose, you may need to use a different form of balancing (a *complete* binary tree is natural for implicit array-based implementations, whereas “left-right height difference cap” is more natural for functional implementations).

1. What do you need to change to transform your *min-heap* implementation (as seen in lecture) to a *max-heap* one?
2. Empirically measure the efficiency of your implementation, like in Lab 4.